

functor category*

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Let \mathcal{C}, \mathcal{D} be categories. Consider the class O of all covariant functors $T : \mathcal{C} \rightarrow \mathcal{D}$, and the class M of all natural transformations $\tau : S \rightarrow T$ for every pair $S, T : \mathcal{C} \rightarrow \mathcal{D}$ of functors. Write $\mathcal{D}^{\mathcal{C}}$ for the pair (O, M) .

For each pair of functors $S, T : \mathcal{C} \rightarrow \mathcal{D}$, write $\text{hom}(S, T)$ the class of all natural transformations from S to T . If τ is in both $\text{hom}(S, T)$ and $\text{hom}(U, V)$, then $S = U$ and $T = V$.

Using the composition of natural transformations, we have a mapping

$$\bullet : \text{hom}(R, S) \times \text{hom}(S, T) \rightarrow \text{hom}(R, T),$$

for every triple $R, S, T : \mathcal{C} \rightarrow \mathcal{D}$. Since composition of natural transformations is associative, the associativity of \bullet applies.

In addition, for each $S : \mathcal{C} \rightarrow \mathcal{D}$, we have the identity natural transformation $1_S \in \text{hom}(S, S)$. For every $\tau \in \text{hom}(S, T)$ and every $\eta \in \text{hom}(T, S)$, we have $\tau \bullet 1_S = \tau$ and $1_S \bullet \eta = \eta$.

From the discussion above, we are ready to call $\mathcal{D}^{\mathcal{C}}$ a category. However, unless $\text{hom}(S, T)$ is a set for every pair of functors in O , $\mathcal{D}^{\mathcal{C}}$ is not a category. When $\mathcal{D}^{\mathcal{C}}$ is a category, we call it the *category of functors* from \mathcal{C} to \mathcal{D} , or simply a *functor category*.

That $\mathcal{D}^{\mathcal{C}}$ is a functor category depends on various restrictions being placed on the “sizes” of \mathcal{C} and \mathcal{D} :

Proposition 1. *If \mathcal{C} is \mathcal{U} -small, then $\mathcal{D}^{\mathcal{C}}$ is a category.*

Proof. Suppose \mathcal{C} is \mathcal{U} -small. Consider the class $\text{hom}(S, T)$. Each $\tau \in \text{hom}(S, T)$ is determined by the collection of morphisms $S(A) \rightarrow T(A)$ for each object A in \mathcal{C} . This means that, for each A in \mathcal{C} , $\text{hom}(S(A), T(A))$ contains the image of every $\tau \in \text{hom}(S, T)$ under A . So the class of all these natural transformations is a subclass of the product

$$\prod_{A \in \text{Ob}(\mathcal{C})} \text{hom}(S(A), T(A)) \tag{1}$$

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Since $\text{Ob}(\mathcal{C})$, as well as each $\text{hom}(S(A), T(A))$ is a set, so is the product (1). Hence $\text{hom}(S, T)$, being a subclass of (1), is a set, or that $\mathcal{D}^{\mathcal{C}}$ is a category. \square

Proposition 2. *If in addition \mathcal{D} is a \mathcal{U} -category, then so is $\mathcal{D}^{\mathcal{C}}$.*

Proof. \mathcal{D} being a \mathcal{U} -category means that $\text{hom}(S(A), T(A))$ is \mathcal{U} -small, for every object A in \mathcal{C} . Since $\text{Ob}(\mathcal{C})$ is also \mathcal{U} -small (assumption in Proposition 1), the product (1) above is \mathcal{U} -small. Consequently, $\text{hom}(S, T)$, being a subclass of (1), is \mathcal{U} -small. This shows that $\mathcal{D}^{\mathcal{C}}$ is a \mathcal{U} -category. \square

Proposition 3. *If \mathcal{D} is furthermore \mathcal{U} -small, so is $\mathcal{D}^{\mathcal{C}}$.*

Proof. We want to show that the class \mathcal{M} of all functors from \mathcal{C} to \mathcal{D} is \mathcal{U} -small. A functor $S : \mathcal{C} \rightarrow \mathcal{D}$ can be broken up into two components: a function $S_1 : \text{Ob}(\mathcal{C}) \rightarrow \text{Ob}(\mathcal{D})$, and a function $S_2 : \text{Mor}(\mathcal{C}) \rightarrow \text{Mor}(\mathcal{D})$, so that $S_2(A \rightarrow B) = S_1(A) \rightarrow S_1(B)$.

Define a binary relation \sim on \mathcal{M} so that $S \sim T$ iff they have the same first component: $S_1 = T_1$. It is easy to see that \sim is an equivalence relation on \mathcal{M} . Let $[S]$ be the equivalence class containing the functor S . For every morphism $A \rightarrow B$, its image under the second component of every functor in $[S]$ lies in $\text{hom}(S_1(A), S_1(B))$. So the size of $[S]$ can not exceed the size of

$$\prod_{A, B \in \text{Ob}(\mathcal{C})} \text{hom}(S_1(A), S_1(B))$$

Since $\text{Ob}(\mathcal{C})$ is \mathcal{U} -small (assumption in Prop 1), so is $\text{Ob}(\mathcal{C}) \times \text{Ob}(\mathcal{C})$. Furthermore, since each $\text{hom}(S_1(A), S_1(B))$ is \mathcal{U} -small (assumption in Prop 2), $[S]$ is \mathcal{U} -small as well.

Next, let us estimate the size of the class \mathcal{M}/\sim of equivalence classes in \mathcal{M} . First, note that for every functor $S : \mathcal{C} \rightarrow \mathcal{D}$, its first component is a function from the set $\text{Ob}(\mathcal{C})$ to the set $\text{Ob}(\mathcal{D})$ by assumption. As $[S] \neq [T]$ iff $S_1 \neq T_1$, the size can not exceed

$$|\text{Ob}(\mathcal{D})^{\text{Ob}(\mathcal{C})}|$$

the cardinality of the set of all functions from $\text{Ob}(\mathcal{C})$ to $\text{Ob}(\mathcal{D})$. By assumption, $\text{Ob}(\mathcal{D})$ is \mathcal{U} -small, so is $\text{Ob}(\mathcal{D})^{\text{Ob}(\mathcal{C})}$. As a result, \mathcal{M}/\sim is \mathcal{U} -small. Together with the fact that $[S]$ is \mathcal{U} -small for each functor S , we have that \mathcal{M} itself must be \mathcal{U} -small, which completes the proof. \square